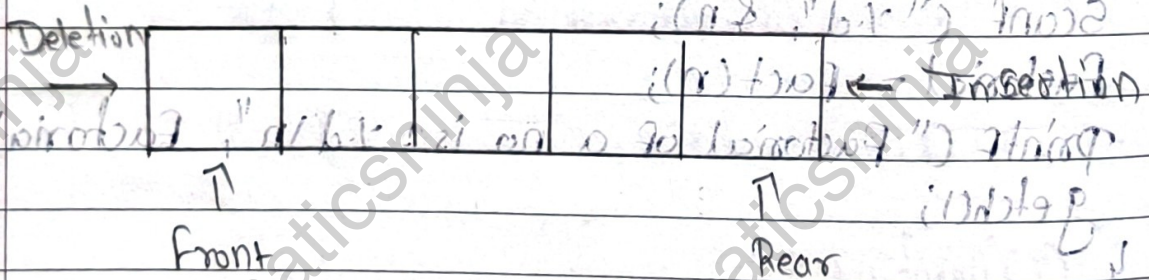


5. Queue

Definition: It is a special kind of list where elements are inserted at one end (rear end) and deleted from the other end (front end).

Queue is the FIFO List :
- FIFO stands for First in First out.
That is the element insert first it will out first.



A Queue datatype may be defined formally as follows:

```
typedef struct queue  
{  
    int data[Max];  
    int front, rear;  
} queue;
```

An array representation of Queue require three entities

- ① An array to hold queue elements.
- ② A Variable to hold index of front element.



③ A variable to hold the index of rear element

During initialization of queue its front and rear are set to -1.

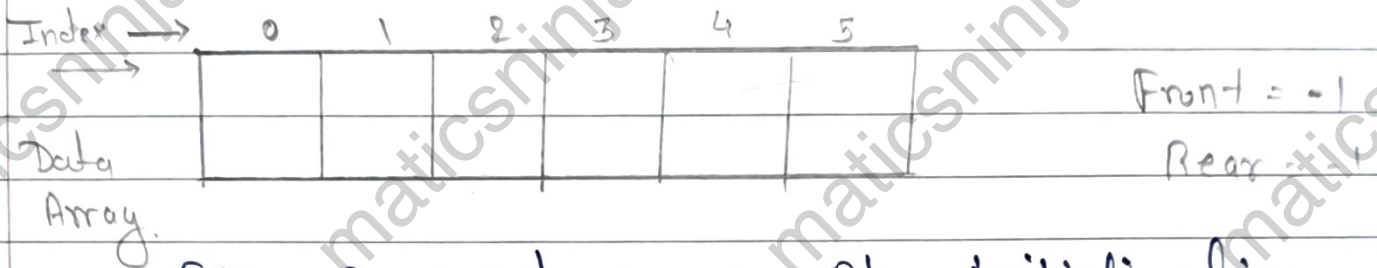
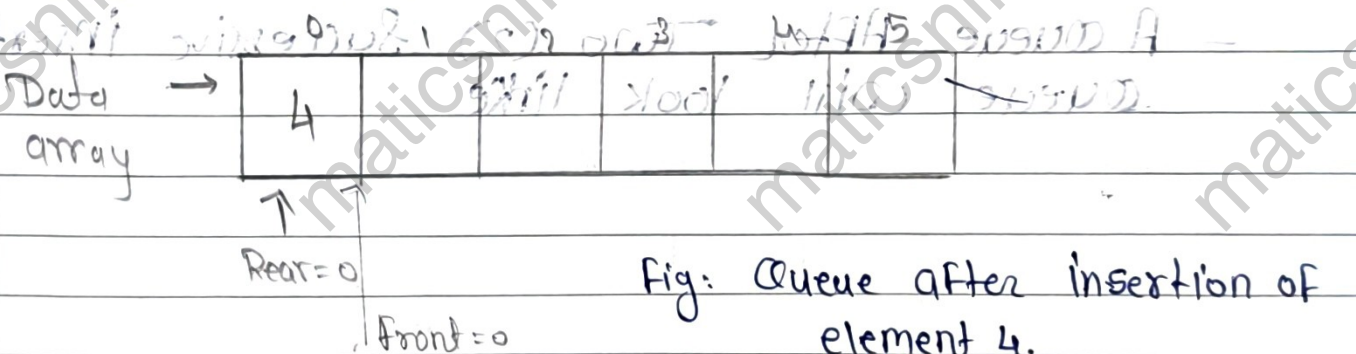


Fig: An empty queue after initialization.

- After insertion of element 4 queue will look like



- After insertion of 3 elements front remains at the same place where rear advances.

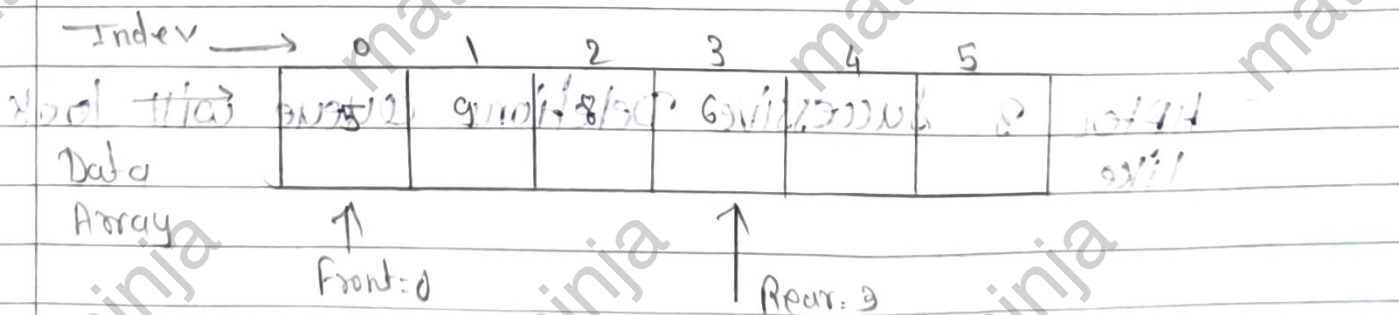


Fig: Queue after insertion of 3 elements is 9, 8, 6.

- A Queue After Three Successive deletions queue will look like

Index	0	1	2	3	4	5	6
Data				6			
Array							

Front = 3 Rear = 3

Fig: A Queue After 3 Successive Deletion.

- A Queue After Two (2) Successive insertion queue will look like

Index	0	1	2	3	4	5
Data				6	5	9
Array						

Front = 3 Rear = 5

Fig: A Queue After 2 Successive Insertion

- After 3 Successive Deletions queue will look like

Index	0	1	2	3	4	5
Data						
Array						

Rear = 1

Fig: A Queue After 3 Successive Deletions.



① if the queue is empty then
 $\text{Front} = 1$ & $\text{rear} = -1$

② if the queue is full
 $\text{rear} = \text{Max} - 1$

③ if $\text{rear} = \text{Front}$ then queue contains just 1
 @ element.

H/W

Q. Distinguish between stack & queue

- Queue overflow & Queue underflow.

Queue overflow

Is caused by insertion in a queue that is already full

Queue underflow

Deletion from an empty queue will cause a queue underflow

enqueue()

enqueue() operation will cause overflow if the queue full.

dequeue()

dequeue() operation will cause underflow if the queue is empty.



operations on queue:

- 1) Initialise()
- 2) enqueue()
- 3) dequeue()
- 4) empty()
- 5) Full()
- 6) print()

1) C function to initialize :

```
Void initialize(Q *p)
```

```
{
    p->rear = -1;
    p->front = -1;
}
```

2) C function to check whether queue full or not.

```
int Full(Q *p)
```

```
{
    if (p->R == max-1)
        return (1);
    return (0);
}
```

3) C function for empty()

```
int empty(Q *p)
```

```
{
    if (p->R == -1)
        return (1);
    return (0);
}
```

→ Distinguish between stack & queue: 1. 10011

Stack	Queue
1] Stack is a LIFO (last in first out) Structure.	1] Queue is a FIFO (first in first out) Structure.
2] In a stack, all insertion & deletion are permitted only at one end of the list.	2] In a queue, items are inserted at one end (the rear) & deleted from the other end front.
3] Stack is always implemented using a linear array.	3] Queue should preferably implemented using a circular array.
4] Stack is widely used in recursion & processing of expression.	4] Queue is more useful in many of the real life applications.
5] It has only one pointer Variable	5] It has two pointer Variable
6] operations: 1. push() 2. pop()	6] operations: 1. enqueue() 2. dequeue()
7] No memory wastage	7] memory wastage in linear queue.

Insert 5 :-

0	1	2	3
5			

Stack

↑ ↑

Rear = 0

Front = 0

↑ ↑

(top of stack)

Stack

Insert 10 :-

Stack is not full. Insertion is possible.

Stack: 5, 10

Front = 0, Rear = 1

↑

Front = 0, Rear = 1

Stack

Insert 7, 8 :-

Stack is not full. Insertion is possible.

Stack: 5, 10, 7, 8

Front = 0, Rear = 3

Stack: 5, 10, 7, 8

Front = 0, Rear = 3

Front = 0, Rear = 3

Delete :-

Stack is not empty. Deletion is possible.

Stack: 10, 7, 8

Front = 1, Rear = 3

Front = 1, Rear = 3

Front = 1, Rear = 3

Stack: 10, 7, 8

Stack



Delete 2 elements

			8
--	--	--	---

Front = 3

rear = 3

1st element = 30

2nd element = 40

3rd element = 50

Delete

--	--	--	--

1st element = 30

2nd element = 40

3rd element = 50

Front = -1

rear = -1

4) C function for enqueue()

Algorithm for insertion in a queue

① Inserting in an empty queue

a) $R = F = 0$

b) $\text{data}[R] = x$

② Inserting in a nonempty queue

$R = R + 1$

$\text{data}[R] = x$

5] dequeue()

Algorithm for deletion from Queue

a) deletion of last element ($R \neq F$)

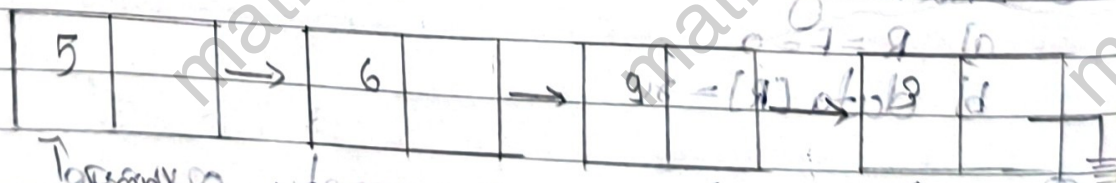
- ① $x = \text{data}[R]$
- ② $R = R - 1$
- ③ $\text{return}(x)$

b) deletion of element where Queue length > 1

- ① $x = \text{data}[F]$
- ② $F = F + 1$
- ③ $\text{return}(x)$

⇒ Array & Link Representation of Queue.

Queue can be represented by Linked Structure of elements. Each element is stored in a node.



• Comparison between Array Representation & Link Representation of Queue.

- Maximum Size of the Queue is fixed at the time of Compilation. When the Queue is Represented using an array this causes wastage of Storage Space.
- In Case of Queue using Linked Structure no memory area is reserved in advanced.
- Memory for node acquire during runtime whenever insertion to be made.
- Array Representation of queue is simple to implement.
- In linked Representation Additional memory required to store address of next element.
- There is no such requirement in Case of array Representation.

→ Types of Queue.

- 1) Linear Queue
- 2) Circular Queue
- 3) Priority Queue.



Advantages of circular queue:

Circular queue is linear data structure in which last node is connected back to the first node to form circle.

It follows FIFO principle.

It is also known as Ring Buffer.

Element are inserted from rear, Deleted from front end.

There is potential problem with implementation of queue using a simple array. The queue may appear to be full although there may be some space in the queue.

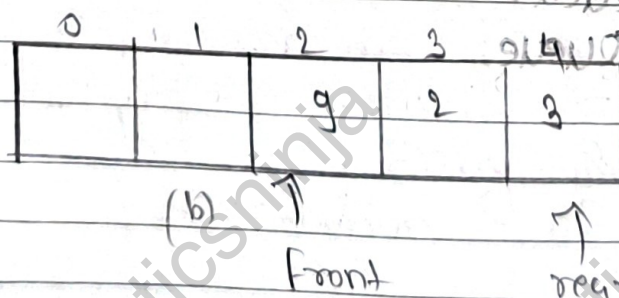
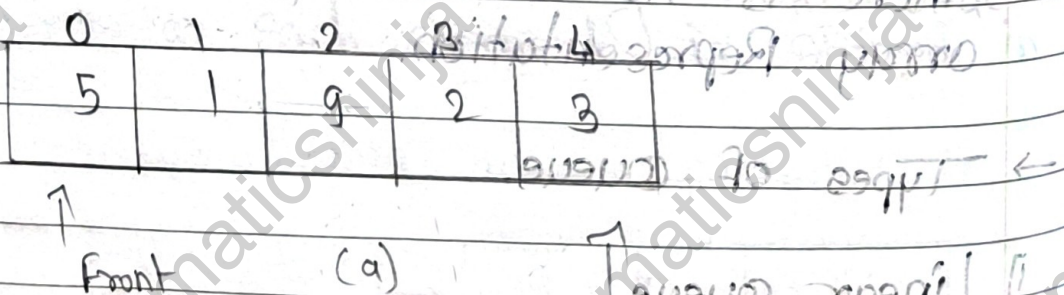


Fig : Queue is full although 2 locations are vacant.



- After insertion of 5 elements in the array is
 $rear = 4$ $front = 0$

- After two successive deletion $front = 2$ and
 $rear = 4$

- Queue in the Fig b is Full no empty space
 head of the rear.

- So To overcome this problem simple solution
 is whenever rear gets to the end of array
 it is wrapped around to the beginning
 Now the array can be thought as a circle

- The first position follows the last element

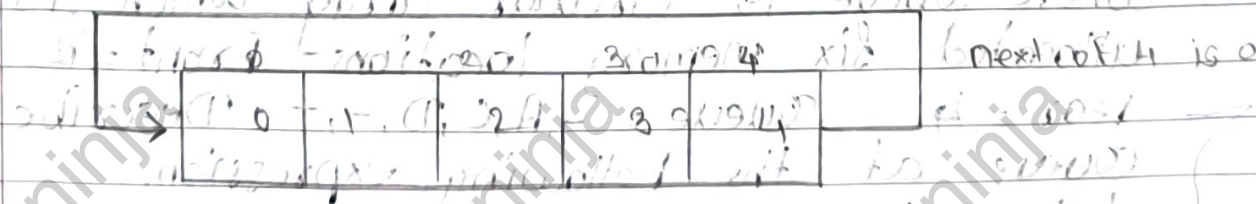
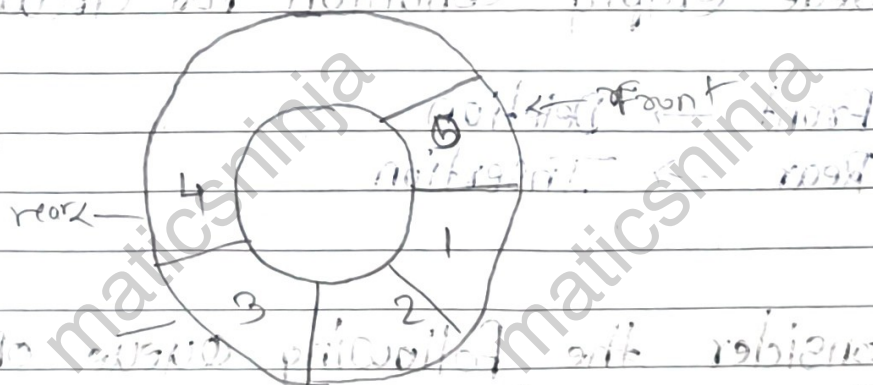


Fig. Circular array.

21. The Condition for Circular Queue Full

$$\text{Front} = (\text{rear} + 1) \% \text{max}$$

0	1	2	3	4
5	4	9	2	8
↑				↑
Front				Rear

$$\text{Front} = (\text{rear} + 1) \% \text{max}$$

$$= (4 + 1) \% 5$$

$$= 5 \% 5$$

$$= 0$$

Q. Explain with Suitable diagram Queue Full and Queue empty Condition for Circular queue.

Front → Deletion
Rear → Insertion

Q. Consider the following Queue of characters where queue is circular array which is allocated six memory locations Front = 2 rear = 4 Queue = -, A, C, D, -, - Describe Queue at the following expression. take place.



operation	Queue	Rear	Front
1] Initially	_, A, C, D, E, _	4	2
2] F is added	_, A, C, D, F, _	5	2
3] Two letters are deleted	_, _, D, F, _	5	4
4] K, L, M are added	L, M, _, D, F, K	2	4
5] Two letters are deleted R is	L, M, _, _, _, K	2	6
6] R is added	L, M, R, _, _, K	3	6
7] Two letters are deleted	_, M, R, _, _, K	3	6
8] one letter is deleted	_, _, R, _, _, K	3	3
9] S is added	_, _, R, S, _, K	4	3

3] Priority Queue

Priority queue is an orderlist of Homogeneous elements. In the normal queue service is provide on the basics of FIFO.

In a priority queue service is not provide on the basics of FCFS, rather than each element has a priority based on the urgency of need.

An element with higher priority is process before all elements with lower priority.

Elements with same priority are processes on FCFS.

Ex:

Hospital waiting queue / Room

Other Application of priority is found in long term scheduling.

• Application of queue.

1] Job / cpu / process scheduling

2] Spooling

3] Queue simulation

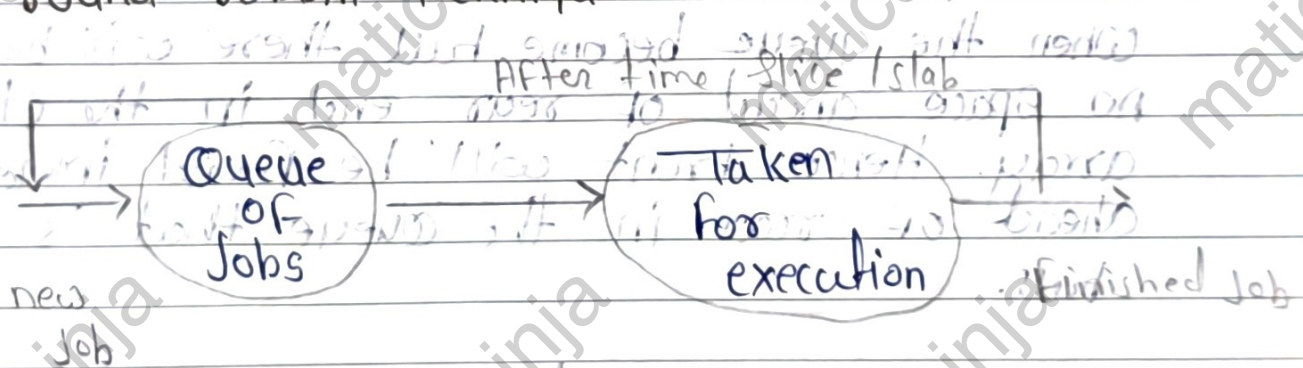
Queue is very useful data structure. Various features of OS (Operating System) are implemented using queue.

1. Scheduling Process (Round Robin)
2. Spooling
3. Queue Simulation

1] Job Scheduling

Jobs entered in the system are stored in queue. The jobs can be selected for execution as per scheduling algorithm.

Round robin technique.



Suppose there are n jobs waiting for execution.

In round robin technique CPU gives fixed Quantum / slice of time to each job.

- ① A job is selected from job queue.
- ② It is executed for fixed slice of time.
- ③ If the job is not completed by this time then it is removed & pushed back in queue of jobs.

Queue Simulation

Queue simulation involves given below

- 1] arrival time
- 2] service time
- 3] Spooling

It is used to maintain queue of job to be printed.

Q. Explain with suitable diagram queue full and queue empty condition for circular queue

→ Queue Full:

When the queue become full there will be no space ahead of rear end in the circular array. Hence front will be found immediately ahead of rear in the queue that is full.

The Queue Full Condition

$$\text{Front} := (\text{rear} + 1) \% \text{max}$$

$$0 = (4+1) \times 5$$

$$Q_{v/p} = 0.5 \times 1.5$$

$$0 = 0$$

5	4	9	2	8
---	---	---	---	---

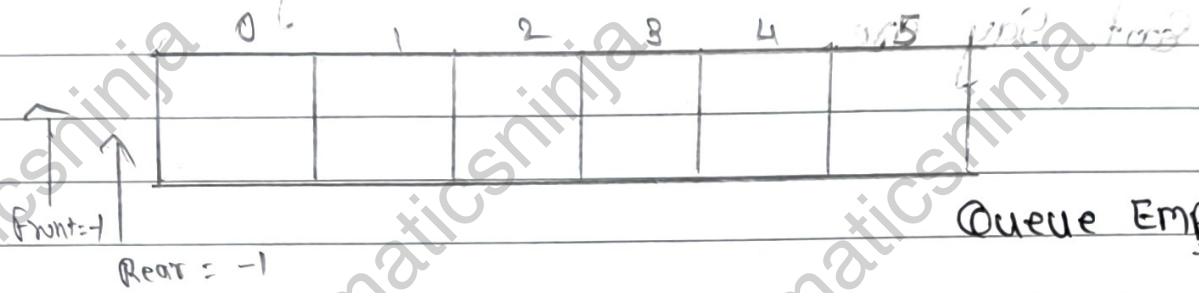
Front

Rear.

Queue Empty

When there is no element found in the queue then queue become empty.

The queue empty condition is $\text{front} = \text{rear} = -1$



When the queue is empty both rear & front will be equal to -1.