

2. Searching and Sorting

→ Searching is the finding location of the data within the given data structure.

- Types of Searching

- I Linear Search

Linear Search is also called as Sequential Search. In Sequential Search all elements are examined sequentially starting from 1st element.

The process of searching terminates when list is ended or comparison results in success.

Algo Search Algorithm starts comparison between $a[i]$ and key value.

- Advantages of Linear Search

- Simple algorithm with less code.
- easy to understand.
- not complex.

Disadvantages of Linear Search

- Should not be used for long list.

Time for searching will increase if list is long.

Best of binary search is based on

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Ex: List = 10, 20, 30, 40, 50 & Search element = 30

Step 1: To search for 30 in 10, 20, 30, 40, 50

10	20	30	40	50
0	1	2	3	4

$10 \neq 30$ Index = Index + 1

Step 2: 10 20 30 40 50

$20 \neq 30$ Index = Index + 1

Step 3: 10 20 30 40 50 → Search is successful
since 30 is found at index 2.

2] Binary Search:

Binary Search is applicable only when the given array is sorted. This method makes comparison between key value that is element to be search and the middle element or centre element of the array.

Since elements are sorted, comparison could be continue with either left half of elements or right half of the elements.

Left half of the elements could be selected by simply making

Right half of the elements could be selected by simply making

process of selecting either the left half or right half continues until the element is found or element is not there.



Ex:

Binary search in array list

Array List: 2, 5, 11, 19, 27, 35, 6, 3, 1, 9

0 1 2 3 4 5

2 5 11 19 27 35 6 3 1 9

Binary Search

i=0

c=2

j=5

 $i = \text{index of first element}$ $j = \text{index of last element}$ $c = \frac{\text{index of first element} + \text{index of last element}}{2}$

$$c = \frac{i+j}{2}$$

$$c = \frac{0+5}{2} = 2$$

2 is from left half

Case 1: if $\text{key} == a[c]$ then element is foundCase 2: if $\text{key} > a[c]$ then right half 19, 27, 35 is selectedCase 3: if $\text{key} < a[c]$ then left half selected

2, 5

from left

selected

2, 5, 11, 19, 27, 35

? ? ? ? ? ?

?

? ? ?

2-1

? - c-1

Q) Search the Element 29 in the given array
 $\{5, 9, 11, 15, 25, 29, 30, 35, 40\}$

$\rightarrow \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{matrix} \rightarrow \text{Index}$
 $\rightarrow \begin{matrix} 5 & 9 & 11 & 15 & 25 & 29 & 30 & 35 & 40 \end{matrix} \rightarrow \text{Element}$

$$i = \text{Index of first element} = 0$$

$$j = \text{Index of last element} = 8$$

$$\text{Centre} = \frac{i+j}{2}$$

$$= \frac{0+8}{2}$$

$$= 4$$

Centre 4 element 25

key Value 29 The element to be
 searched is key Value equal to 29

3. If Step 2d: if $\text{key} > a[c]$ then $i = c+1$
 $29 > 25$ then $i = 6$

The right half part is selected

29, 30, 35, 40	Element
5 6 7 8	Index
↑ ↑ ↑	
i=5 c=6	J=8

$$C = \frac{i+j}{2} = \frac{5+8}{2} = \frac{13}{2} = 6 \text{ which is } \\ \text{key value at index}$$

Step 2: Key is less than $a[8]$, that is $29 < 30$.

Left half part will be selected.

$29 \rightarrow \text{Element}, 5 = \text{Index}$

∴ left half part contain only one element
that is key value

∴ Search is successful.

∴ key value is get on location 5 that is 29

Q. Find the position of element 29 using Binary search in array given below

$$A = \{11, 5, 21, 3, 29, 17, 2, 43\}$$

$i < j$

→ Sorted list of given arry:

Data should be sorted for binary search.

$$\begin{array}{ccccccccc} 2 & 3 & 5 & 11 & 17 & 21 & 29 & 43 \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \end{array} \rightarrow \text{Element}$$

$i = 0$

$j = \text{Index of last element} = 7$

$c = \frac{0+7}{2} = \text{Algorithm to find } c$

Centre = 3

element = 11

key value is 29. The element to be search in right part.

Step 1: Key > a[c] i.e. 11
 $29 > 11$

So the left half is discarded.

The right half part is selected.

Elements are also written below block 4721.

8 4 5 6 7 → Index

17 21 29 43 → Element

↓
C

↓

So left & right are the part of sorted part.

$$c = \underline{4+7} = 5$$

original array 29 43 17 6 7 11 8 3 2 5 4 → sorted array 29 43 17 6 7 11 8 3 2 5 4

Step 2: Search movie position of 29

Key > 21

Key > a[c] i.e. 8, 10, 3, 11 & 1

$$29 > 21$$

Left part is selected.

Elements of left part 11 8 3 10

29 43 17 6 7 11 8 3 10

6 7

$$\therefore c = \underline{6+7} = \underline{13} = 6 \text{ what is } 1$$

∴ from 2nd block in what is 1

∴ The (element) position of element 29 is found at position 6

∴ 29 is found at position 6

⇒ Sorting = ~~Arranging~~ ~~values~~ ~~tail~~ ~~to~~ ~~order~~
 Sorting is the Arranging data in a logical order.

- Types of Sorting.

- i) Bubble Sort

Bubble Sort is one of the simplest and most popular sorting method.

The basic idea behind bubble sort is as a bubble rises up in the water, the smaller element goes to the beginning.

This method is based on successive selecting element through exchange of adjacent element.

Let n be the number of elements in array the 1^{st} pass begins with the comparison of a_1 of $n-1$, and a_2 of $n-2$. If a_2 is larger than a_1 of $n-1$ the two elements are exchanged.

Ex: P-1: 5 2 1

To ~~arrange~~ ^{arrange} millions fast we do it
 Swap two said if ne forming

P-2: 2 1 5



1 2 5

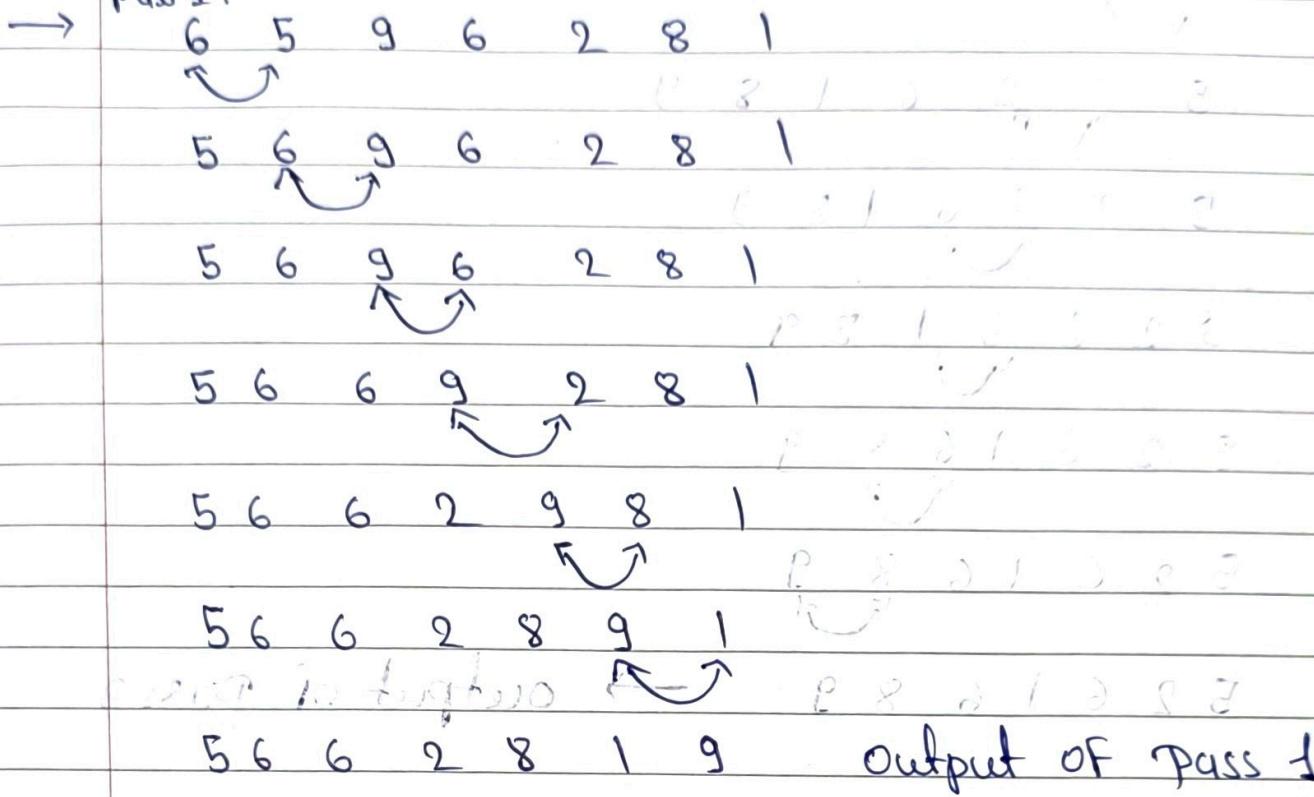
1 2 5

arrange in ~~arrange~~

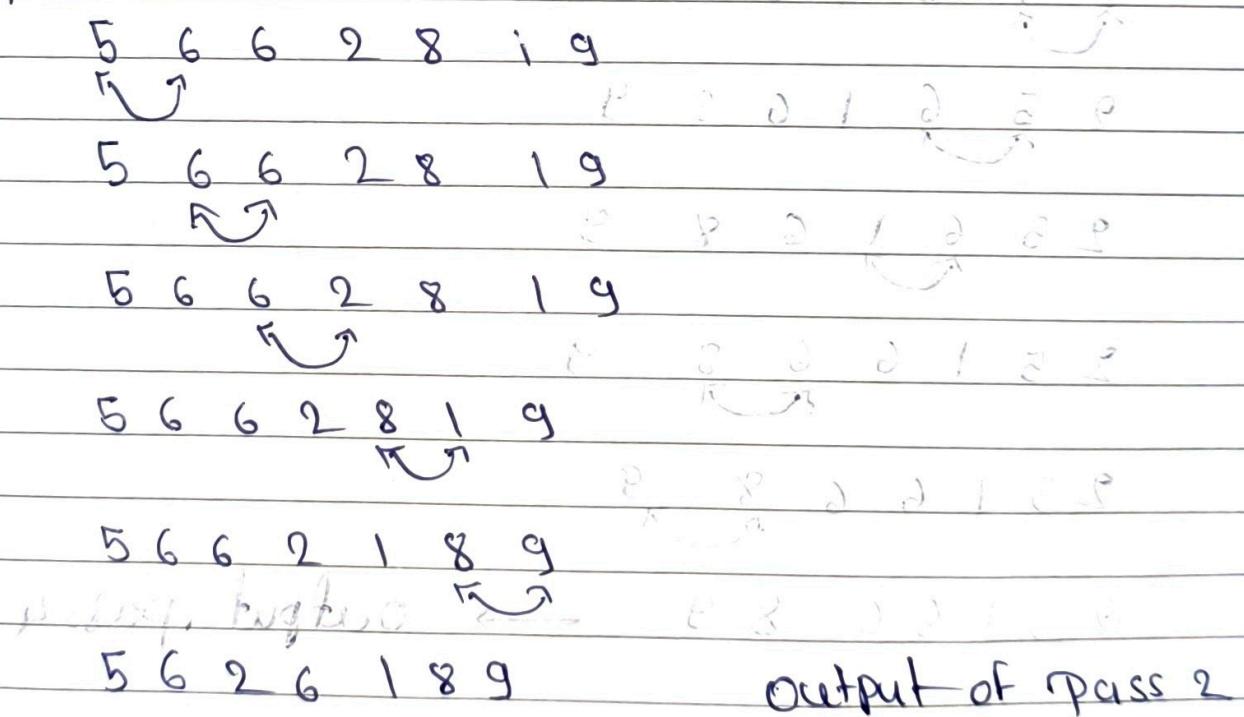


Q. Original array with $n = 6, 5, 9, 6, 2, 8, 1$

Pass 1:



Pass 2:



Pass 3: 3.2.B.2.0 in New place Torripico

5 6 2 6 1 8 9

1 3 8 9 0 2 8 3

5 6 2 6 1 8 9

1 3 8 9 0 2 8 3

5 2 6 6 1 8 9

1 3 8 9 0 2 8 3

5 2 6 6 1 8 9

1 3 8 9 0 2 8 3

5 2 6 1 6 8 9

1 3 8 9 0 2 8 3

5 2 6 1 6 8 9

1 3 8 9 0 2 8 3

5 2 6 1 6 8 9

→ output of pass 3

Swap to first

Pass 4:

5 2 6 1 6 8 9

P 1 2 0 2 3 7

2 5 6 1 6 8 9

P 1 2 0 2 3 7

2 5 6 1 6 8 9

P 1 2 0 2 3 7

2 5 1 6 6 8 9

P 1 2 0 2 3 7

2 5 1 6 6 8 9

P 2 1 0 2 3 7

2 5 1 6 6 8 9

→ output pass 4

Swap to first



Pass 5: in another partition with head 16
from oldest price to new

2 5 1 6 6 8 9 1 8 10 9 11 3 11

in 2 5 1 16 6 8 9 in another partition with head 16
from oldest price to new

2 1 5 11 6 6 8 9 10 9 11 2 11 3 11

2 1 5 6 6 8 9

2 1 5 6 6 8 9 10 9 11 2 11

2 1 5 6 6 8 9

10 9 11 8 10 9 11 2 11

Pass 6:

2 1 5 6 6 8 9 10 9 11 2

1 2 5 6 6 8 9 = Sorted List

1 2 5 6 6 8 9 10 9 11 2

Swapping logic:

$\text{temp} = a[j]$

$a[j] = a[j+1]$

$a[j+1] = \text{temp}$

1 2 5 6 6 8 9 10 9 11 2

1 2 5 6 6 8 9 10 9 11 2

1 2 5 6 6 8 9 10 9 11 2



2] Selection Sorting

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Selection Sorting is the Very Simple Sorting method.

In the i^{th} pass we will select the element with lowest value among $a[i], a[i+1], \dots, a[n-1]$ and swap it with $a[i]$. As a result after i passes elements will be sorted order.

Ex: 8210 810 278 811 21 21

Original Array 5 9 1 11 2 4

5 9 1 11 2 4

After pass 1 \rightarrow 1 9 5 11 2 4

After pass 2 \rightarrow 1 2 5 11 9 4

After pass 3 \rightarrow 1 2 4 11 9 5

After pass 4 \rightarrow 1 2 4 5 9 11

After pass 5 \rightarrow 1 2 4 5 9 11

Q1) Find no. in ascending order

2. Sort the no in ascending order
 $A = \{42, 23, 74, 11, 65\}$

→ original array 42 23 74 11 65

The Small no is 11
 Comparing with $a[0]$ and exchange it.

After Pass 1: 11 23 74 42 65

The Small no is 42
 Comparing with $a[1]$ & exchange it.

After Pass 2: 11 23 42 74 65

The Small no is 65
 Comparing with $a[3]$ & exchange it.

After Pass 3: 11 23 42 65 74

→ Sorted array

11 23 42 65 74 →

37

Insertion Sort

An Element can always be placed at ~~right~~ place in the Sorted List.

Ex: 5 9 10 15 20

Element to be placed ~~is 6~~ after 10
if the element 6 is to be inserted in sorted list of elements its ~~right~~ place between 5 and 9

Elements with values greater than 6 should be moved right by 1 place.
Thus creating a space for incoming element.

Insertion Sort

Q. Arrange the given no. in Ascending order using insertion sort.

77 33 44 11 88 22 66 55

→ [77 33 44 11 88 22 66 55] → Initial unsorted array

Sorted Unsorted

Next element is 33 insert before 77

in sorted array

[33 77] [44 11 88 22 66 55] → After Pass 1

Sorted

Unsorted

Next element is 44 insert between 33 & 77

[33 44 77] [11 88 22 66 55] → After pass 2

Sorted

Unsorted

Next element is 11 insert before 33

in sorted array

[11 33 44 77] [88 22 66 55] → After Pass 3

Sorted Unsorted

Next element is 88 insert before 77 in
sorted array

[11 33 44 77 88] [22 66 55] → After Pass 4

Next element is 22 insert between
11 & 33 in sorted array

11 22 33 44 77 88 | 66 55 → After Pass 5
 Sorted Unsorted

Next element is 66 insert in between 44 & 77 in
 Sorted array.

11 22 33 44 66 77 88 | 55 → After Pass 6
 Sorted Unsorted

Next element is 55 insert in between 44 & 66 in
 Sorted array.

11 22 33 44 55 66 77 88 |
 0P 1P 2P 3P 4P 5P → After Pass 7
Sorted list

Q1) Radix Sort : -

Radix Sort is generalization of bucket sort to sort decimal numbers where the radix or base is 10. we need 10 bucket

In the first pass numbers are sorted on ~~list~~ least significant digit. numbers with the same least significant digit are stored in same bucket in the.

In the 2nd pass numbers are sorted on 2nd least significant digit.

at the end of every pass numbers in buckets are merge to produce a common list

Number of passes required to sort is equal to number of digits in the largest number in the list.

Q. Sort the given no in ascending order using quick sort.

Sort. 348, 14, 614, 5881, 47

→ Pass 1:

Merge list : 5381, 14, 614, 47, 348; 1 : fail 39339

Pass 2:

Merge List: 14, 614, 47, 348, 5881 / 0

Pass 3:

47	5381
14	348
0	614

14, 49, 348, 5381, 614

merge list :

14, 47, 348, 5381, 614

Pass 4:

the sorted list will be formed by merging with first
 614, 848, 147, 348, 5381

147, 348, 5381

0	1	2	3	4	5	6	7	8	9
---	---	---	---	---	---	---	---	---	---

mergelist : 14, 47, 348, 614, 5381.

This is the sorted list formed by merging with first

and process will go on till all elements are merged.

(b) forming sorted list by V to R pass

V to R pass

S1 S2 S3 S4 S5 S6

first merging + S1 S2 S3 S4 S5 S6

(2) forming sorted list

S1 S2 S3 S4 S5 S6

last step of merging

S1 S2 S3 S4 S5 S6

S1 S2 S3 S4 S5 S6

5] Quick Sort:

Quick Sort is the internal Sorting Algorithm with the time complexity equal to $O(n \log n)$

The basic algorithm to sort an array of n elements can be described recursively as follow

Step 1: if $n <= 1$ return

Step 2: Pick any element v in array $a[]$
this is called the pivot element

Step 3: rearrange elements of the array by moving all elements x, v

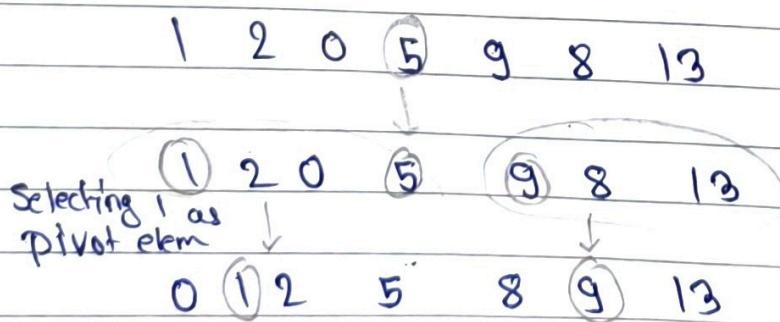
~~Step 4:~~ Right of v & all the elements x, v
left of v

①. 5, 1, 2, 9, 8, 13

→

5 1 2 9 0 8 13 ← original list
array

Select Pivot element ⑤



This is sorted list by numbers 2, 5, 8, 9, 13
Using quick sort and binary search
algorithm to search it in sorted

Time complexities of different algorithm

1] Binary Search.

Time complexity: $\log_2 n$

2] Sequential Search

Time complexity: $O(n)$

3] Bubble Sort: To compare last with first
Time complexity: $O(n^2)$

4] Insertion Sort

Time complexity: $O(n^2)$

5] Selection Sort to remove off in bubble

Time complexity: $O(n \log n)$

6] Quick Sort: Divide & Conquer and merge

Time complexity: $O(n^2)$